

# *Collective Efficiency in Two-sided Matching*

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# The Highlight of the Talk

Markets such as labor market, a venue for bilateral trading, require a proper matching between potential traders.

- (1) Propose new matching algorithms based on the **bounded rationality**.
- (2) Investigate the scaling-property of the matching models.
- (3) Application in the trading model of the Sugarscape world.

# A Model of Two-sided Matching

## <The marriage model>

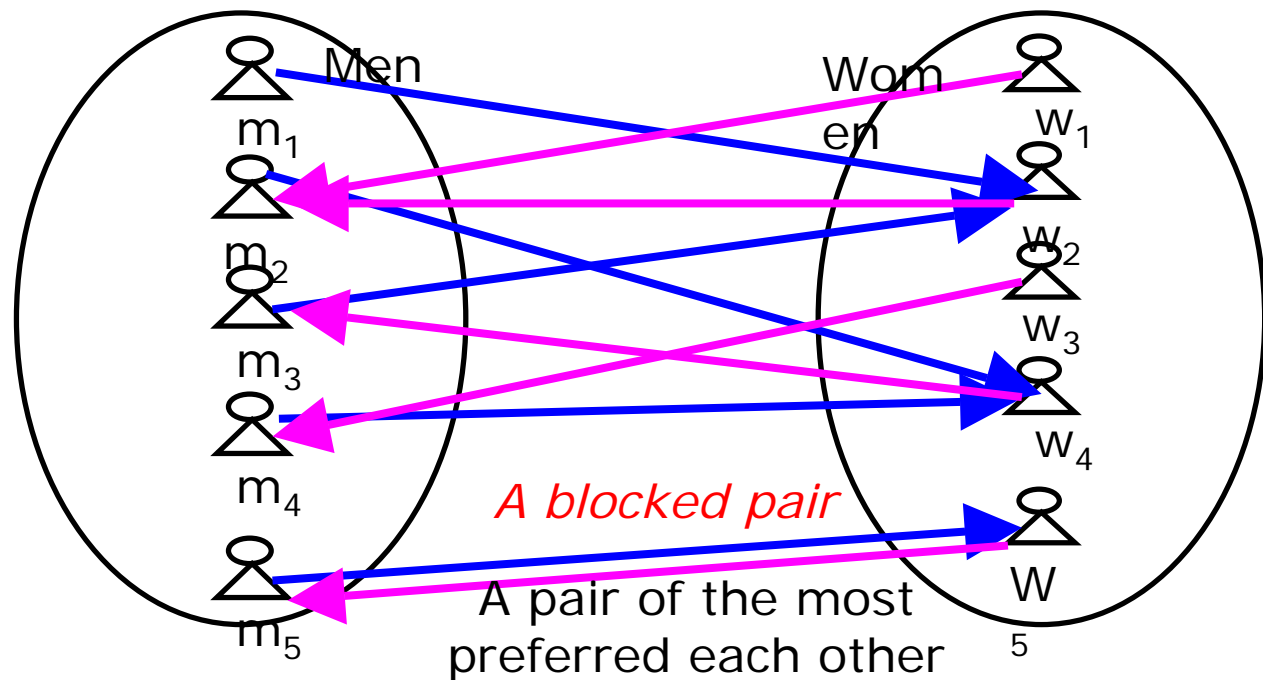
$$G_A(\text{Men}) = \{m_1, \dots, m_n\}, G_B(\text{Women}) = \{w_1, \dots, w_p\}$$

- Preference order of a man ( $m_i$ ) over  $G_B$

$$P(m_i) = w_3, w_2, \dots, w_j, \dots$$

- Preference order of a woman ( $w_j$ ) over  $G_A$

$$P(w_j) = m_2, m_4, \dots, m_i, \dots$$



# Stable Matching

*A matching  $\tau$  is blocked* by a pair of agents  $(m,w)$

if they each prefer each other to their mates at some different matching  $\eta$  .

A matching  $\tau$  *is STABLE* if it isn't blocked by any pair of agents.

<Stable matching> **(Nash equilibrium)**

Equilibrium of **individual interest-seeking behaviors**

# Deferred Acceptance Algorithm (DA)

<Step 1>

- a. Each man  $m$  proposes to his 1st choice.
- b. Each woman who receives more than one proposes, she "holds" the most preferred one.

<Step  $k$ >

- a. Any man who was rejected at step  $k-1$  makes a new proposal to his most preferred to acceptable mate who has not yet rejected him.
- b. Each woman holds her most preferred acceptable offer to date, and rejects the rest.

# Properties of DA

**<Theorem>** (Gale and Shapley 1968)

The matching DA produces is always stable with respect to the strict preferences.

**<Theorem (Knuth)>**

The best outcome for one side of the market is the worst for the other.

- (1) **Man-optimal** stable matching is **the worst for women**.
- (2) **Women-optimal** stable matching is **the worst for men**.

# Scaling Properties of DA

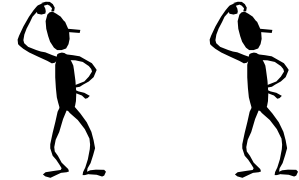
(1) The population sizes:  $N$  (*Omero, Dzierzawa:2004*)

(2) Agent  $i$  (man) and agent  $j$  (woman) are paired.

$x_i$ : rank order of agent  $j$  in the women

$y_j$ : rank order of agent  $i$  in the men

$$1 \leq x_i, y_j \leq N$$



$x(y)=1$  is the best and  $x(y)=N$  is the worst

(3) Definition of energy (level of being unsatisfied)

$$\text{Energy of men } E_M = \sum x_i \quad \text{Energy of women } E_F = \sum y_j$$

(4) Scaling property based on random preferences

$$E_M \cong \log N \quad E_F \cong N / \log N \quad E_{total} = E_M \times E_F \cong N$$

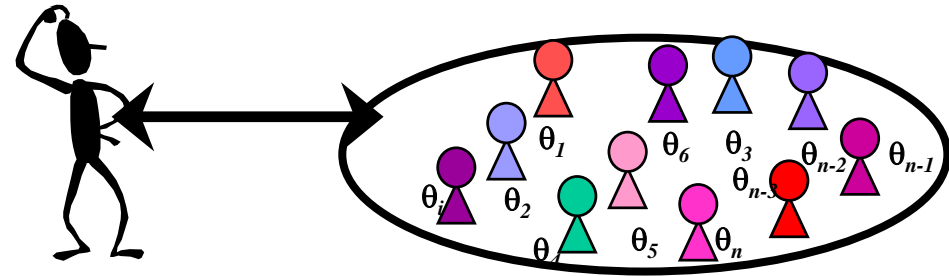
(5) The total energy of DA and the minimum energy

$$F_{total} = E_M + E_F$$

$$F_{total}^{stable} = 2\sqrt{N} \quad F_{total}^{min} = 1.62\sqrt{N}$$

*The performance of DA is about 80% of the optimal matching*

# Matching Based on Bounded Rationality



**(1) Model 1: Accept if the offer is within some compromise level**

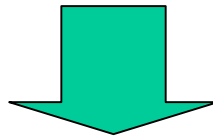
**(2) Model 2: Reject the offer if it below than some threshold**

**(1) No Compromise or no threshold**

**Deferred Acceptance Algorithm**

**(2) Too much compromise (or low threshold)**

**Accept any offer (or hold any offer)**



**The optimal compromise level or threshold is set so that the social utility (collective efficiency) is maximized**



# ASP [Acceptance with Some Patience]

<Step 1>

- a. Each man proposes to his 1st choice.
- b. Proposed woman accepts the proposal within **the acceptance level**, otherwise rejects it. **(Holding is not allowed)**

<Step k>

- a. Each not matched woman increases the acceptable level by one
- b. Any man who was rejected at step k-1 makes a new proposal to his most preferred acceptable mate who has not yet rejected him.
- c. Proposed woman accepts the proposal within the acceptable level, otherwise she rejects it.

*Acceptance level = Compromise level + Step number*

# Definitions of Utilities

N: The population size

$x_j$ : rank order of agent j in the women

$y_i$ : rank order of agent i in the men

$x(y)=1$  is the best and  $x(y)=N$  is the worst

(1) Utility of a paired individuals i and j

$$U_i = 1 - y_i / N \quad U_j = 1 - x_i / N$$

(2) Utility of a pair of individuals i and j

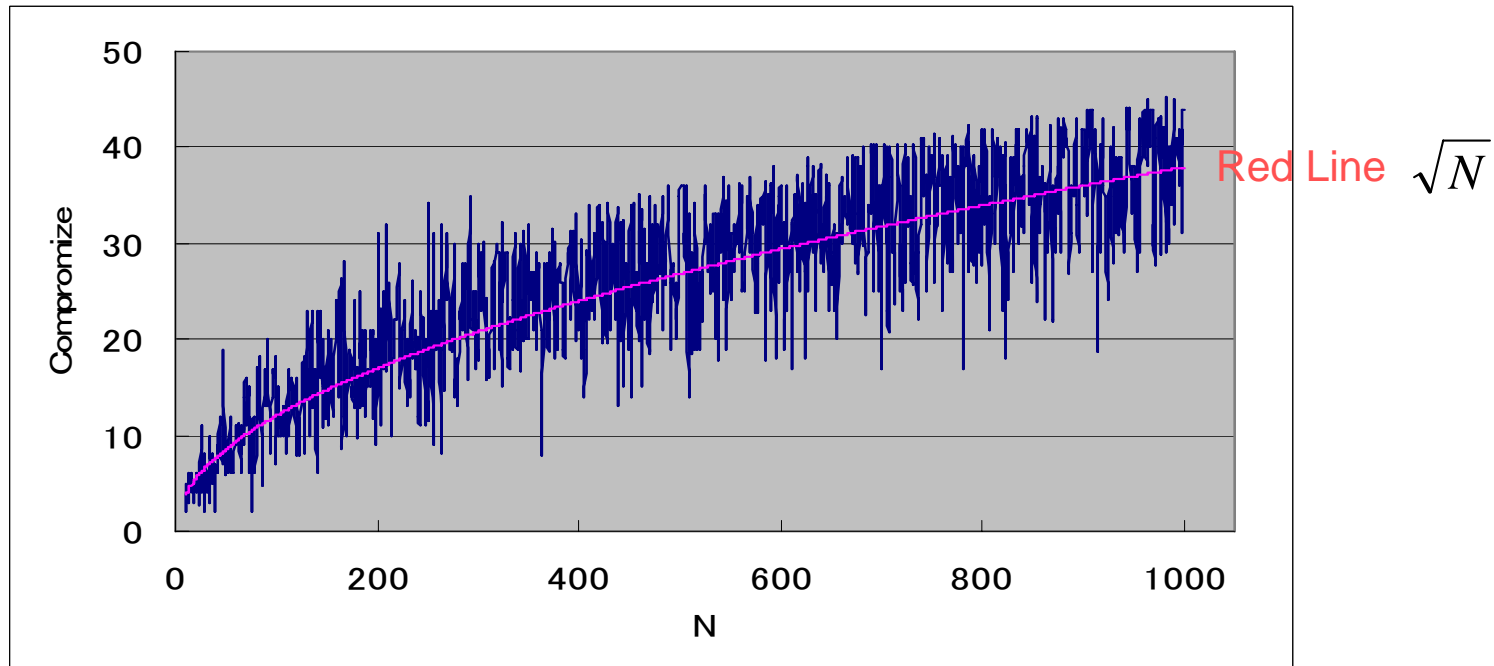
$$U(p_{ij}) = U_i \times U_j$$

(3) Social Utility (Collective welfare)

$$SU = \sum_{p_{ij}} U(p_{ij})$$

# The Optimal Compromise Level

Blue Line: The compromise level  $C$  to maximize social utility



<The Square Root Law of Compromise>

The optimal compromise level  $\lambda(N)$ :  $\lambda(N) = \sqrt{N}$

# Advanced Deferred Algorithm (ADA)

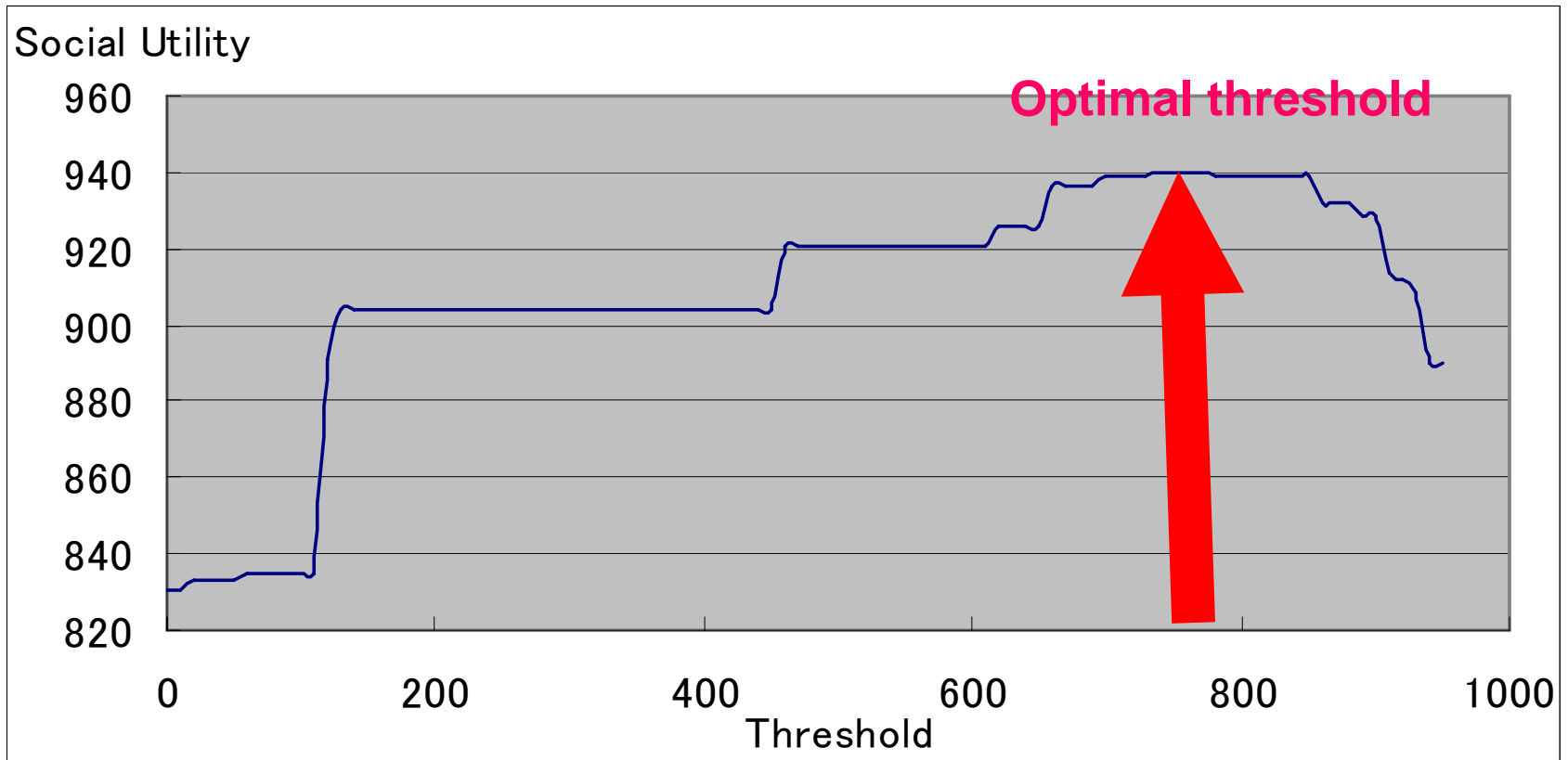
<Step 1>

- a. Each man  $m$  proposes to his 1st choice.
- b. Each woman “rejects” the propose if it is lower than some “threshold” , otherwise “hold” the most preferred one.

<Step  $k$ >

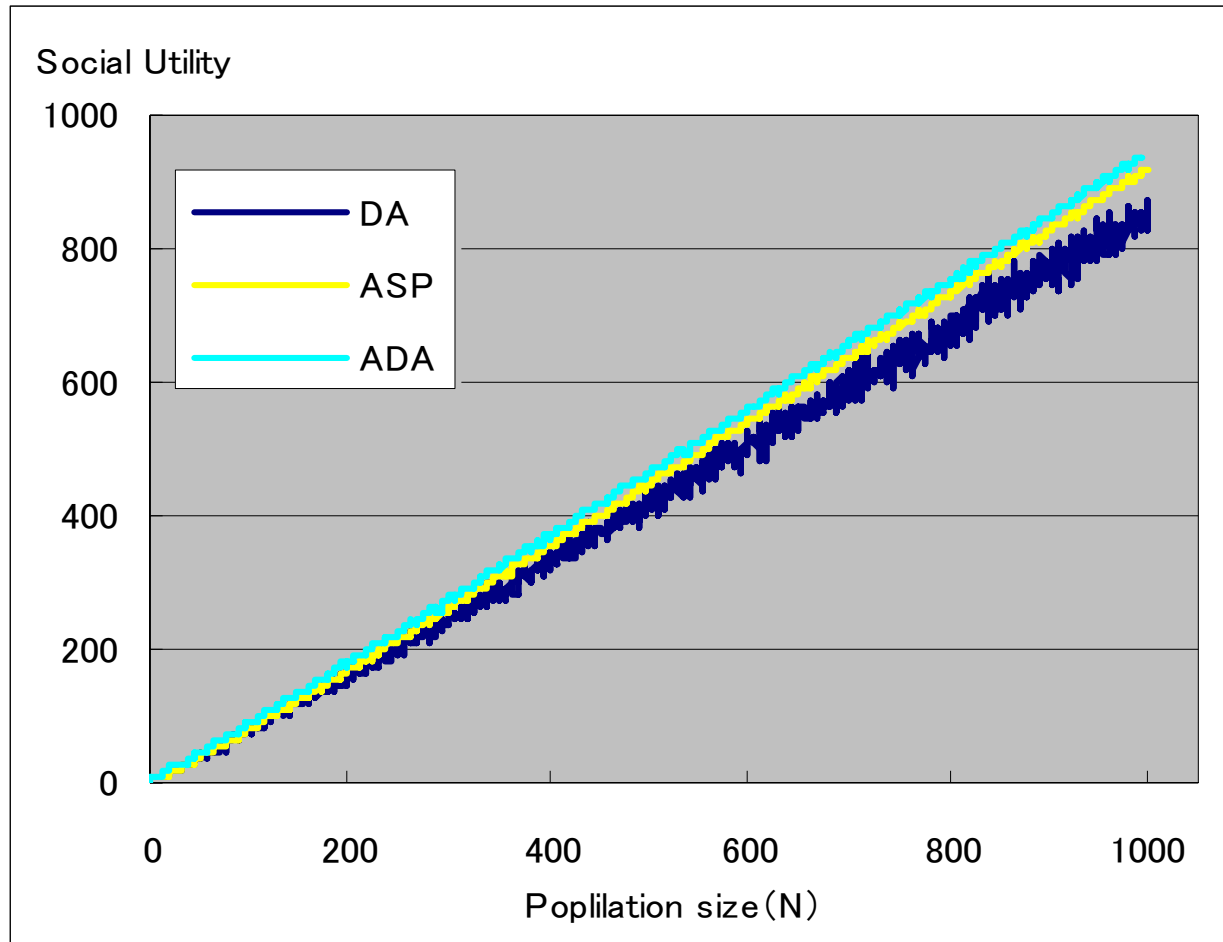
- a. Any man who was rejected at step  $k-1$  makes a new proposal to his most preferred to acceptable mate who has not yet rejected him.
- b. Each woman holds her most preferred acceptable offer to date, and rejects the rest.

# The Optimal Threshold of the Cutoff



(The population of 1,000 agents)

# Comparison of Collective Efficiency

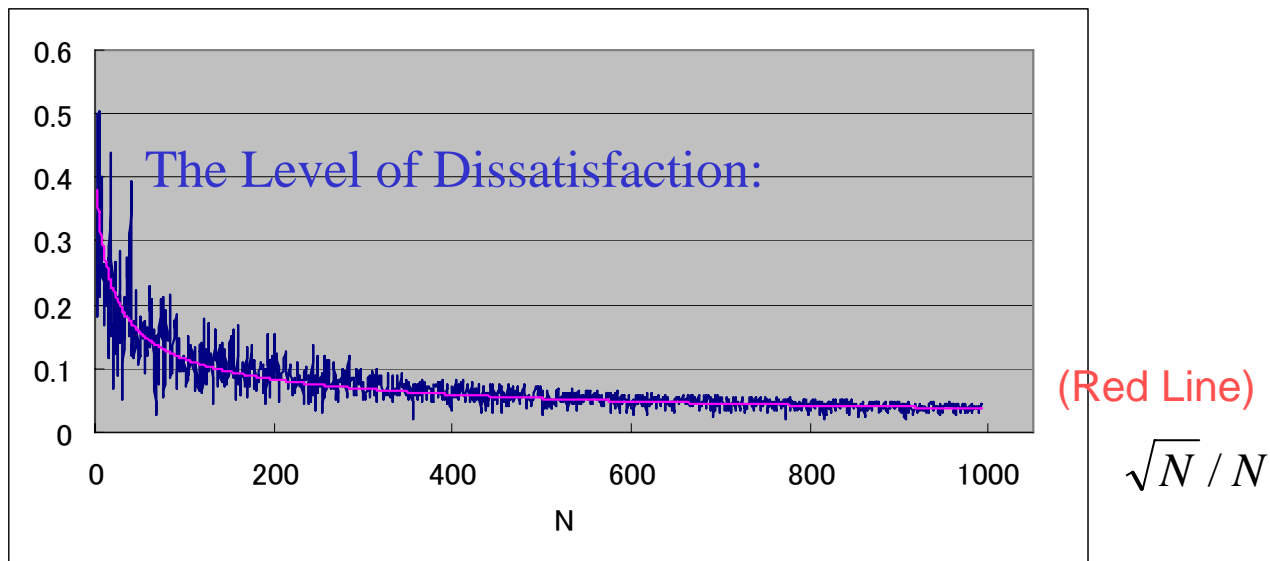


The maximum efficiency: 1

- (1) Two-sided matching with compromise (ASP): 0.92
- (2) Advanced Deferred Acceptance (ADA): 0.93
- (3) Deferred algorithm (DA): 0.85

# The Scaling Property of Collective Efficiency

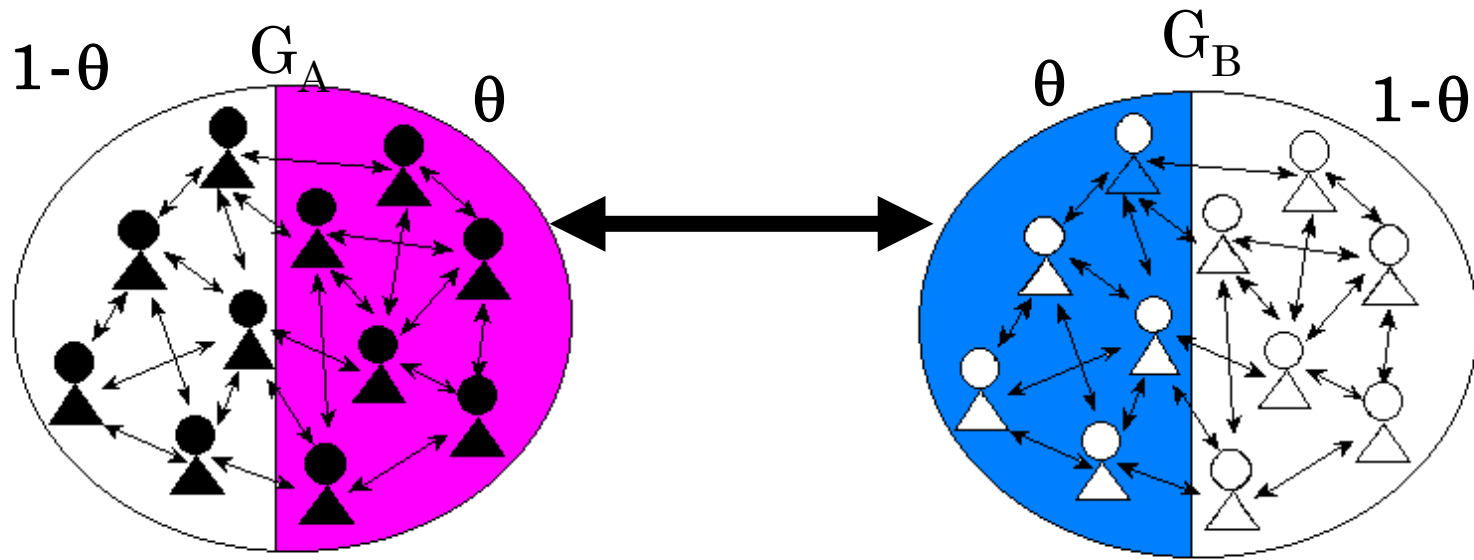
$$\text{Agent average utility} \cong 1 - (\sqrt{N} / N)$$



*Individually rational behavior results in collective rationality if the population size increase!!*

*Then why are there so many mismatches in our societies?*

# Random Preferences vs. Correlated Preferences



<Random preferences >

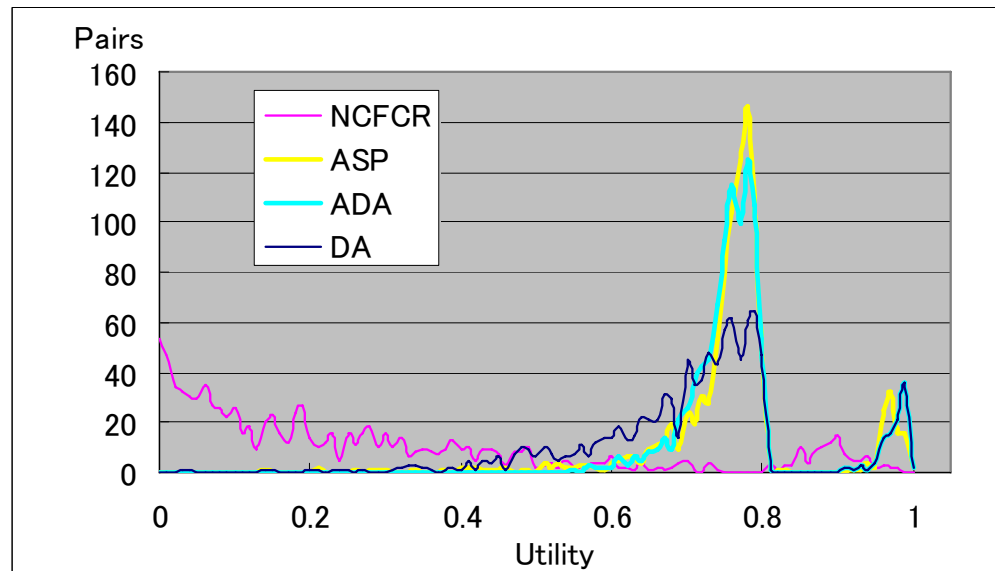
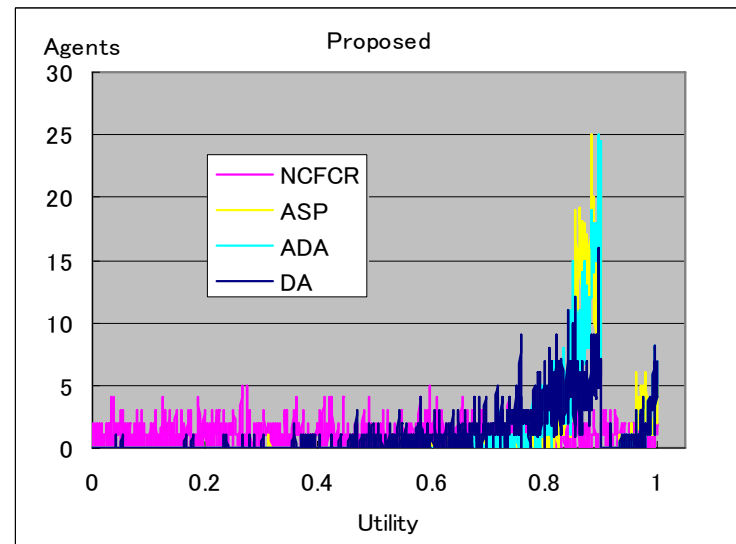
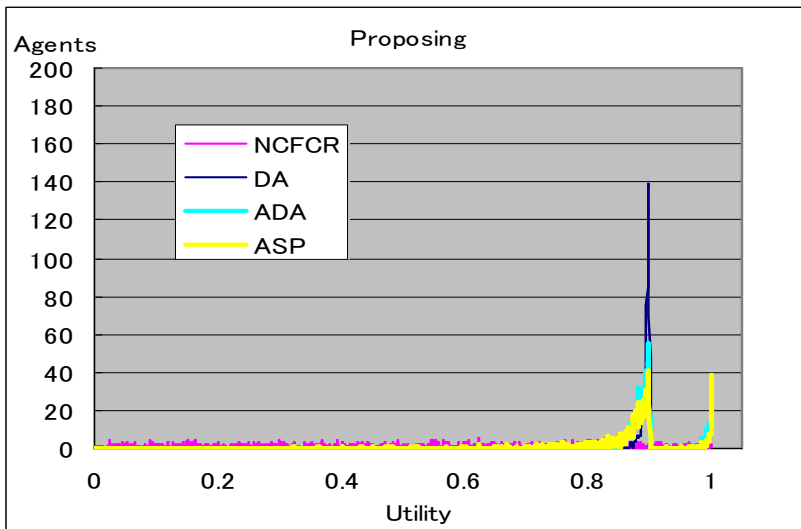
(1) Preferences of men and women are generated randomly.

<Correlated preferences >

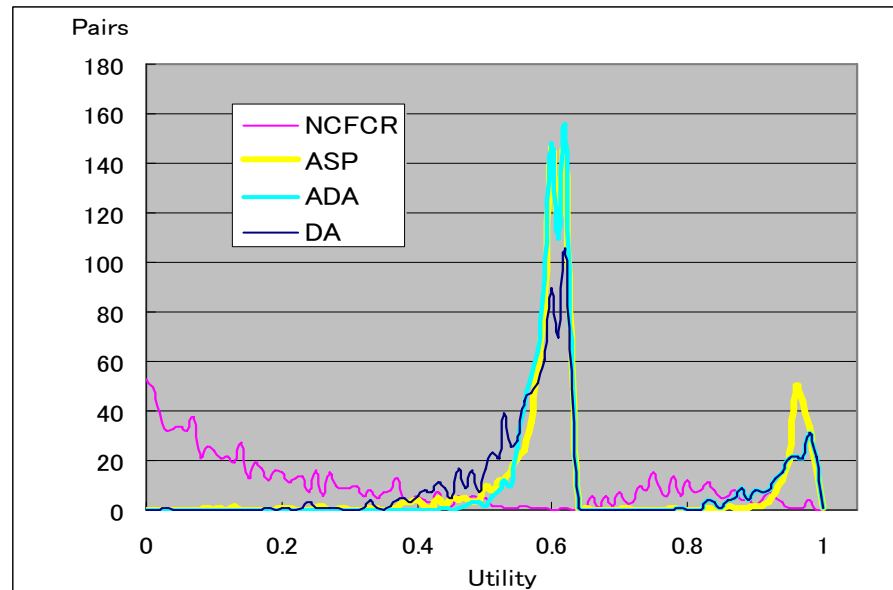
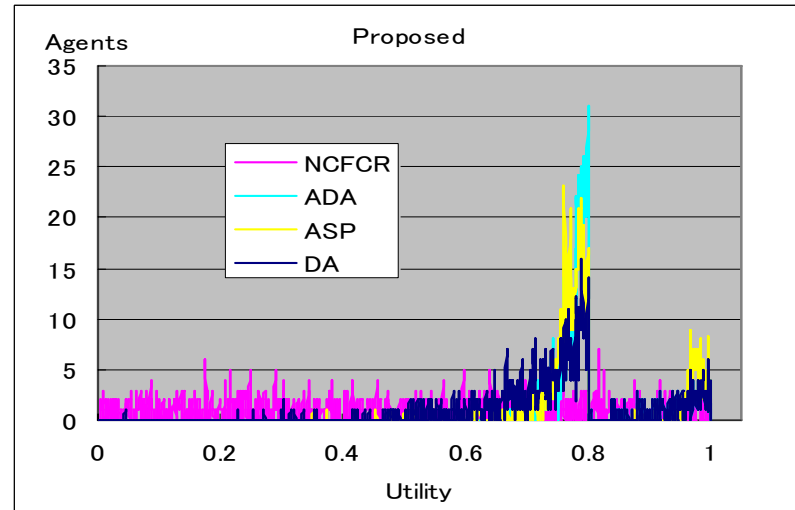
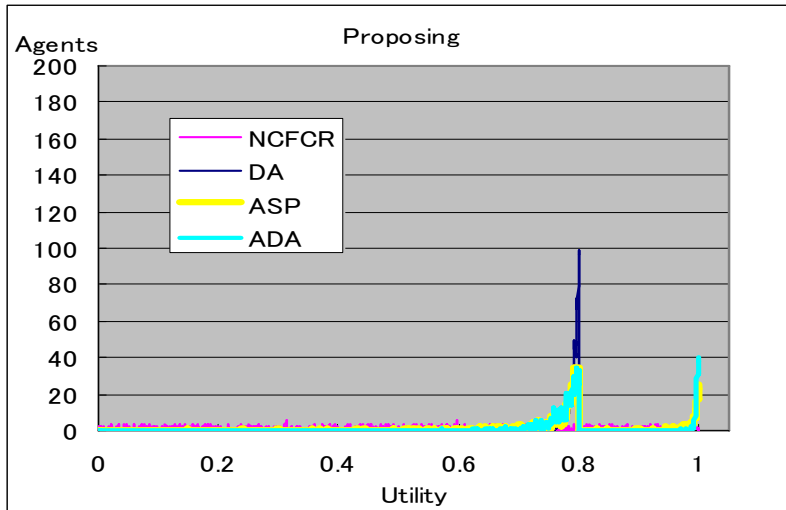
- (1) There are top stars (both men and women) who are preferred by all members of the other side. (The ratio of top stars:  $\theta$ )
- (2) The rest of the preferences (the ratio of  $1-\theta$ ) are generated randomly



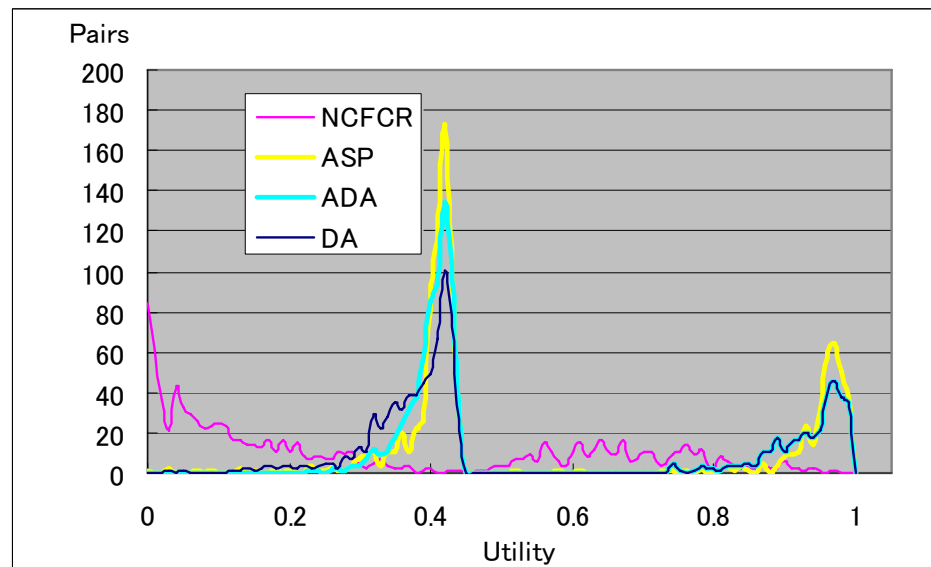
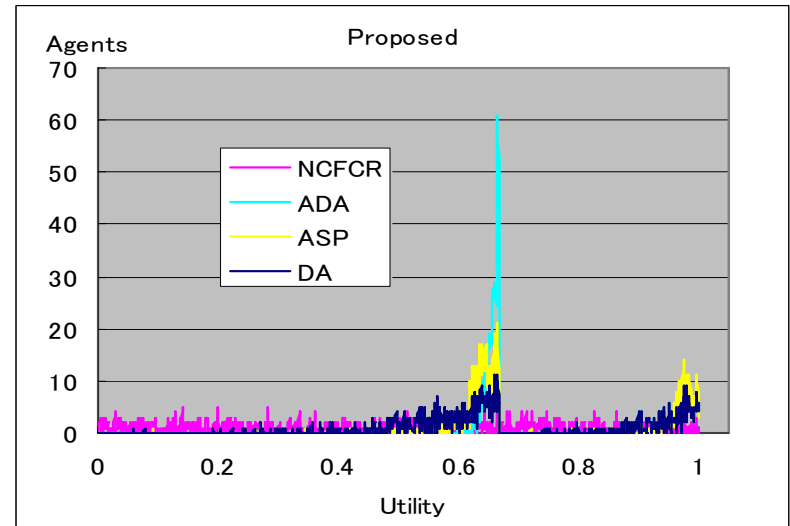
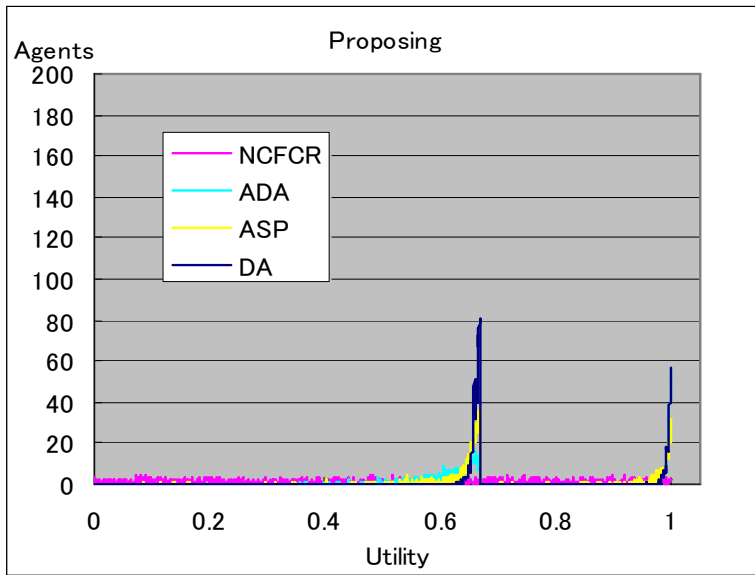
# The Results under Correlated Preferences : Case 1 ( $\theta = 0.1$ )



# Correlated Preferences: Case 2 ( $\theta = 0.2$ )

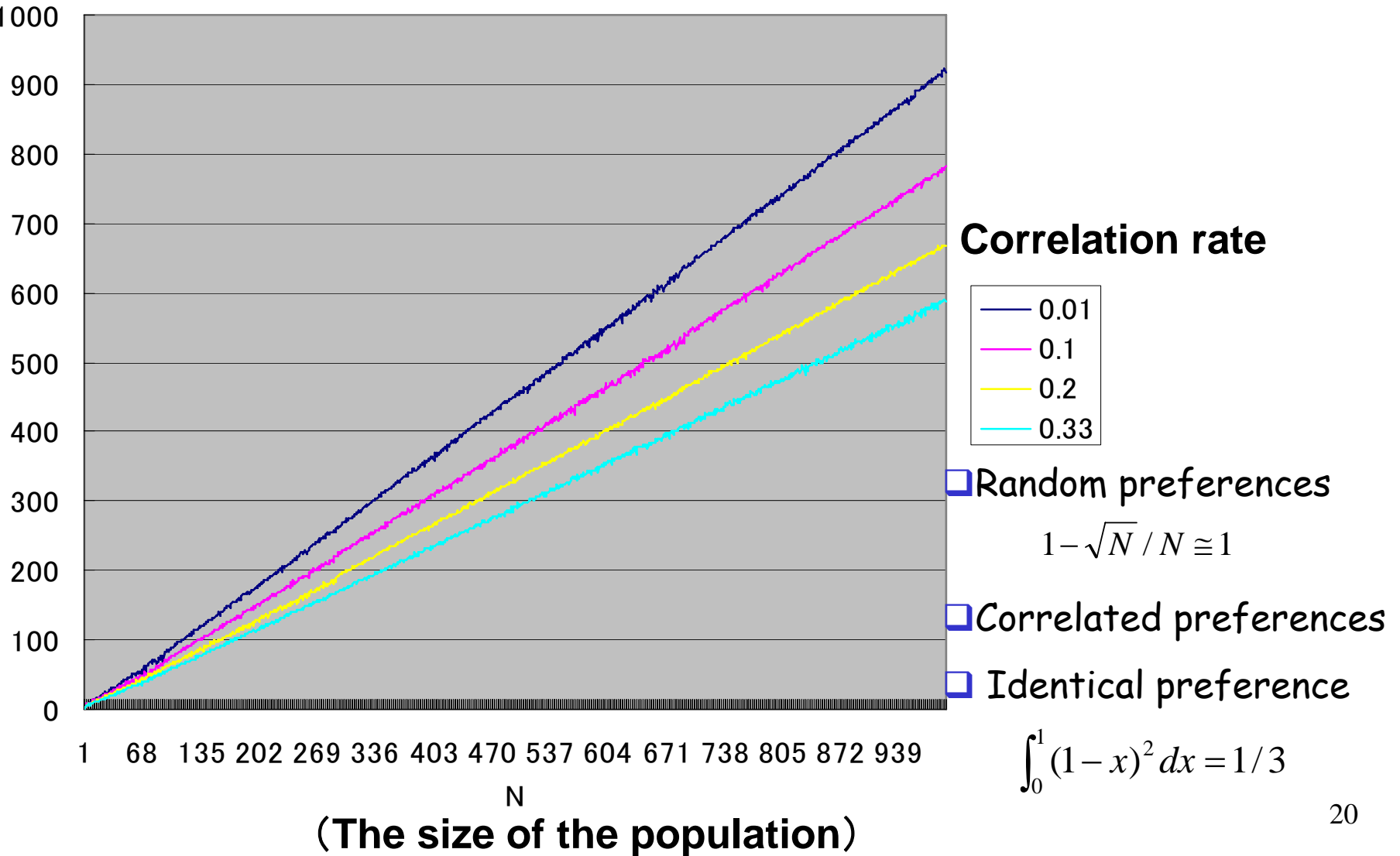


# Correlated Preferences: Case 3 ( $\theta = 0.33$ )



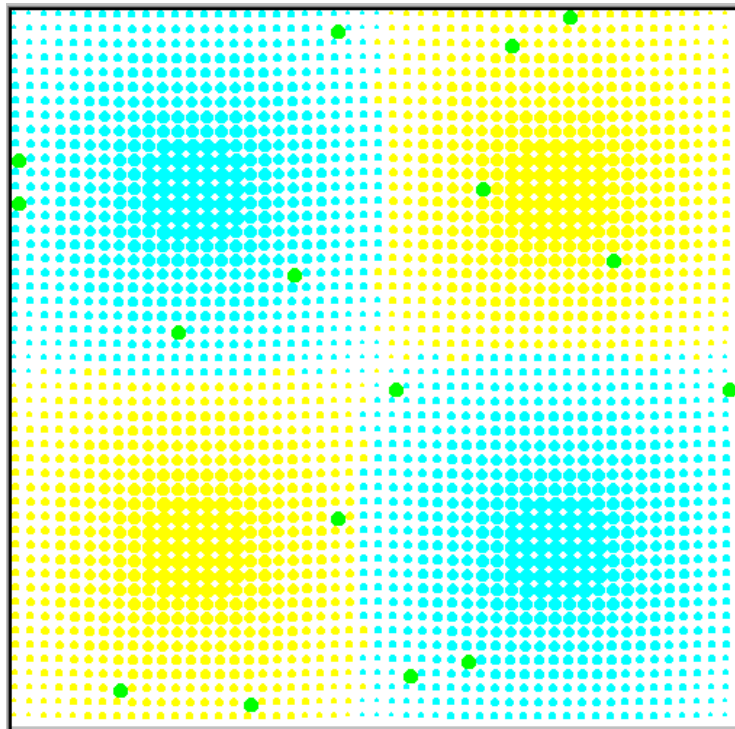
# Summary:




## Collective Efficiency under Correlated Preferences



# Application of Two-sided matching in the Artificial Society

M. Epstein and R. Axtell developed the Sugarscape model for the study of new social science from bottom-up (1996)



Agent   
Resource A   
Resource B 

They need two resources A and B to survive.

(Epstein & Axtell Model)

**Epstein & R. Axtell showed that equilibrium price is achieved through local trading without market.**

# Bilateral Trading in the Sugarscape model

Agents trade resources to survive



Metabolism for resource:

Stock of resource:

Resource A    Resource B



$m_A, m_B$

$W_A, W_B$

Utility function in terms of resources A and B :

$$W(w_A, w_B) = w_A^{m_A/m_r} w_B^{m_B/m_r} = w_A^a w_B^{1-a}$$

The Marginal Rate of Substitute:  
(The relative price of resource A to B)

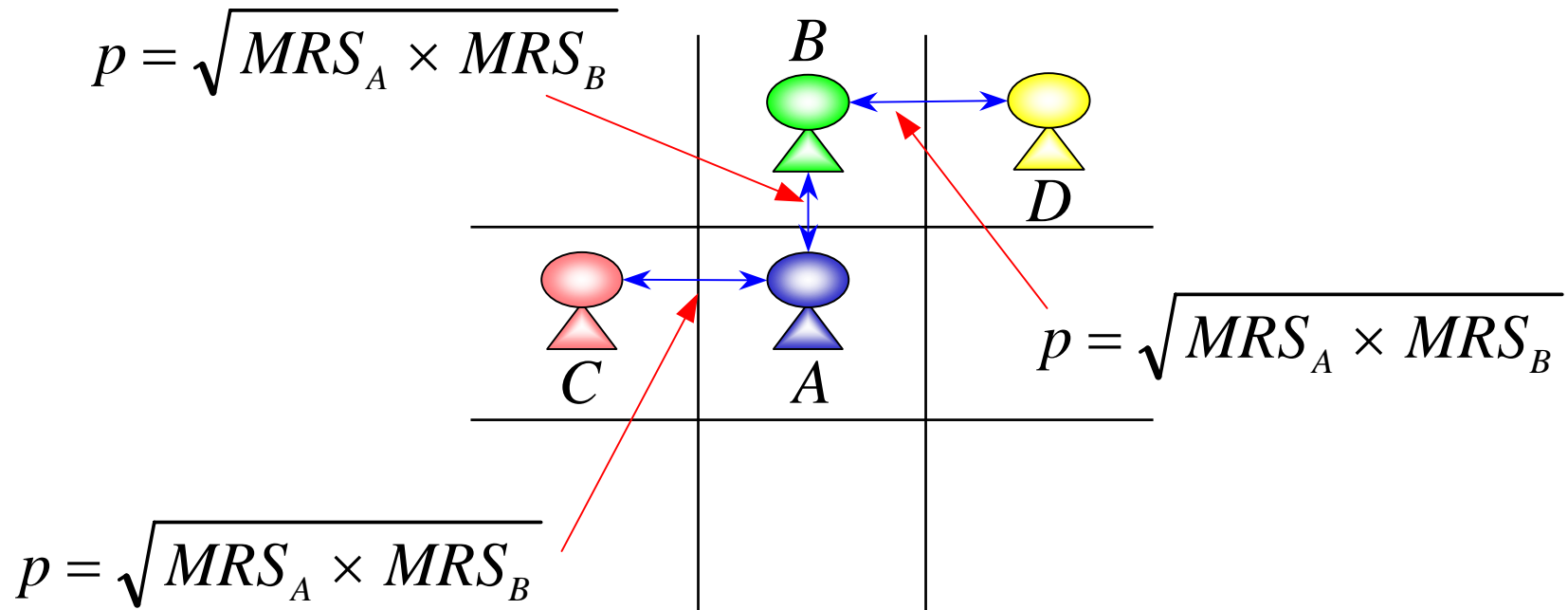
$$MRS \equiv \frac{\partial U / \partial w_A}{\partial U / \partial w_B} = \frac{w_B / m_B}{w_A / m_A}$$

$MRS > 1 \rightarrow$  An agent needs resource A

$MRS < 1 \rightarrow$  An agent needs resource B

# Bilateral Trading: Local Pricing

## Model 1: Epstein & Axtell Model



The price is determined by the geometric average of  $MRS$  of two agents

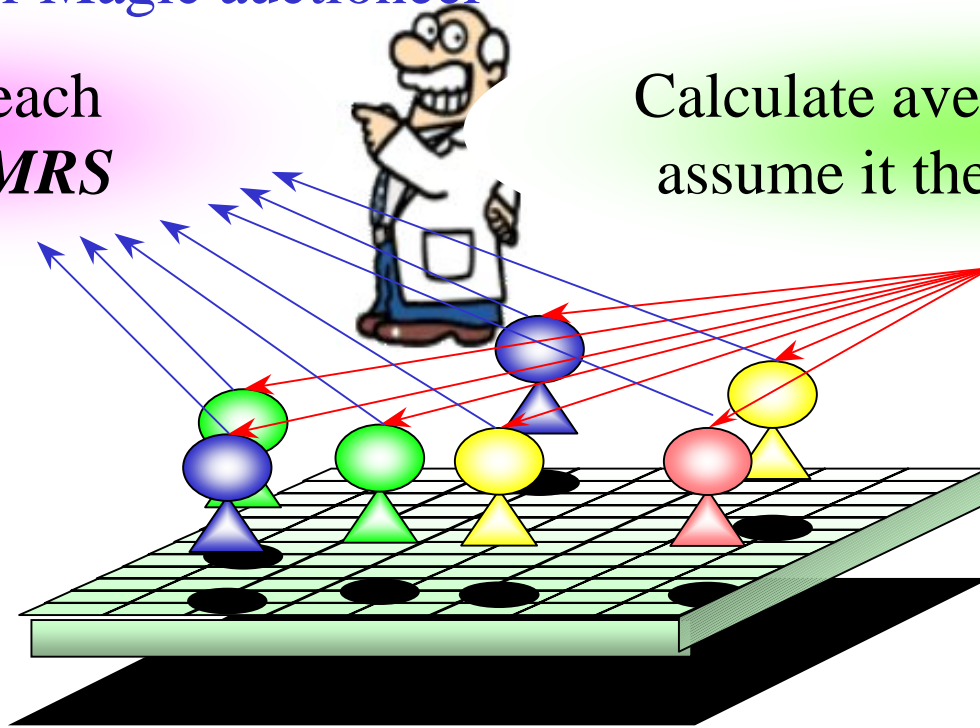
$$p = \sqrt{MRS_A \times MRS_B}$$

# Global Pricing Model (Model 2)

Global market or Magic auctioneer

Collect each agent's *MRS*

Calculate average *MRS* and assume it the market price



The price is decided by averaging *MRS* of all agents

$$p = \sum_i^N MRS / N$$



# Bilateral Trading by Two-sided Matching (Model 3)

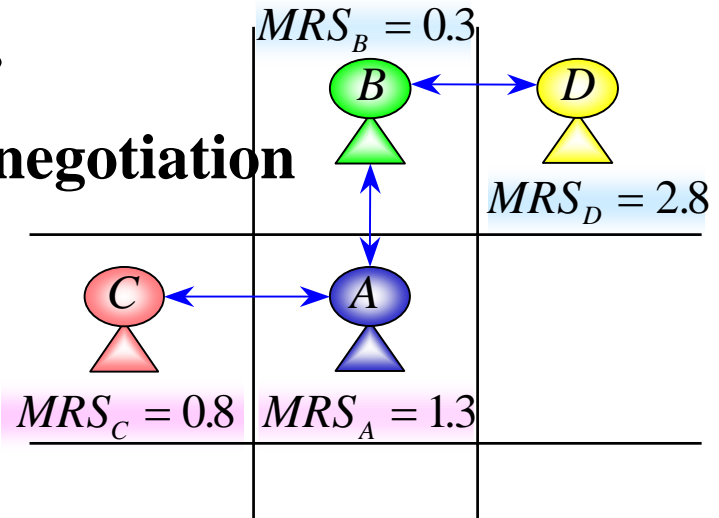
(1) Trading between agent  $i$  and agent  $j$  will take place if

$$MRS_i > 1 \text{ and } MRS_j < 1.$$

(2) Trading will take place under negotiation

$S_1$ : demand partner's  $MR$

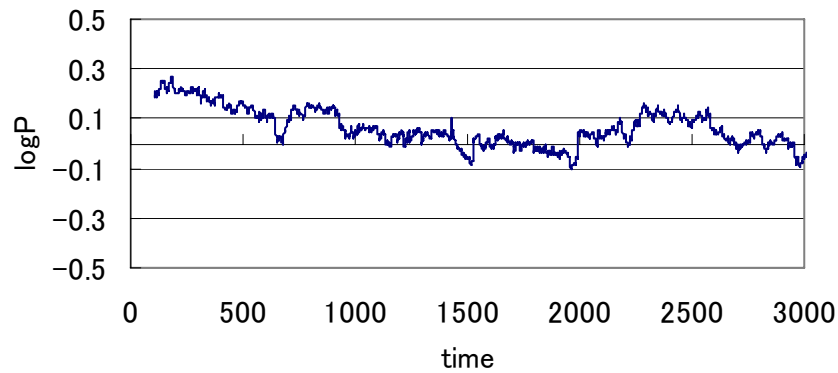
$S_2$ : forgo with own  $MRS$



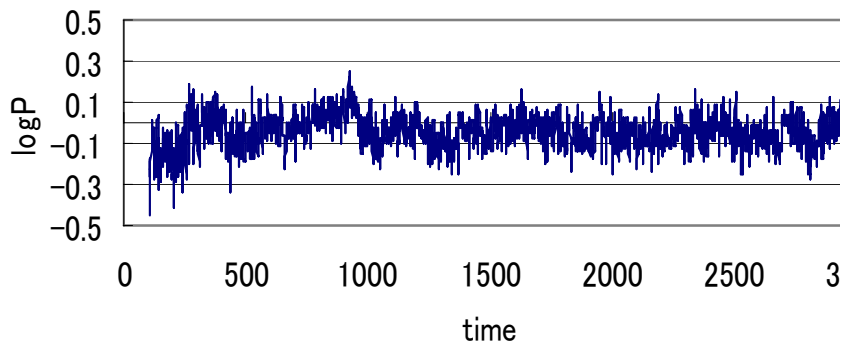
		Agent $j: MRS < 1$	
		Hawk	Dove
Agent $I: MRS > 1$		$S_1$	$S_2$
Hawk	$S_1$	Geometric average of $MRS$	demand partner's $MRS$
Dove	$S_2$	Own's $MRS$	geometric average of $MRS$

# Comparison of Trading Prices

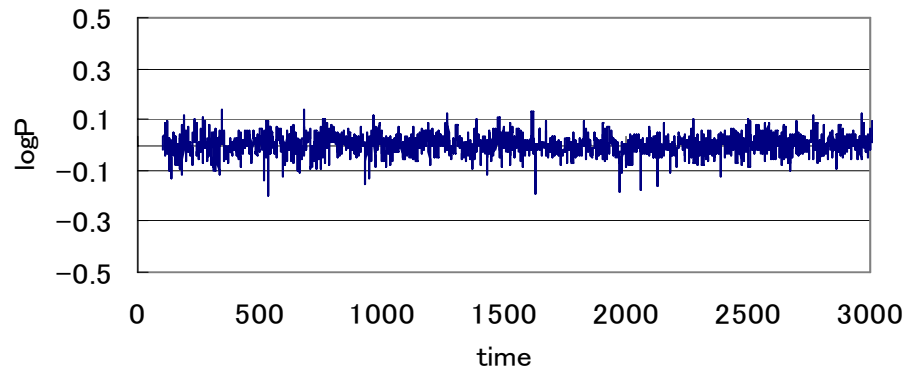
Model 2: Global pricing



Model 1: Bilateral trading

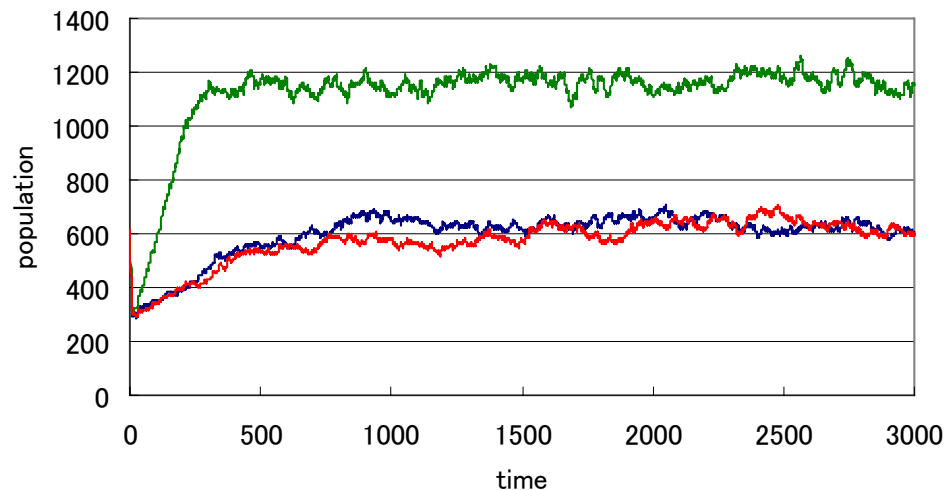


Model 3: Bilateral trading with two-sided matching



# The Population Size and Average Utility (Efficiency)

<Population size>



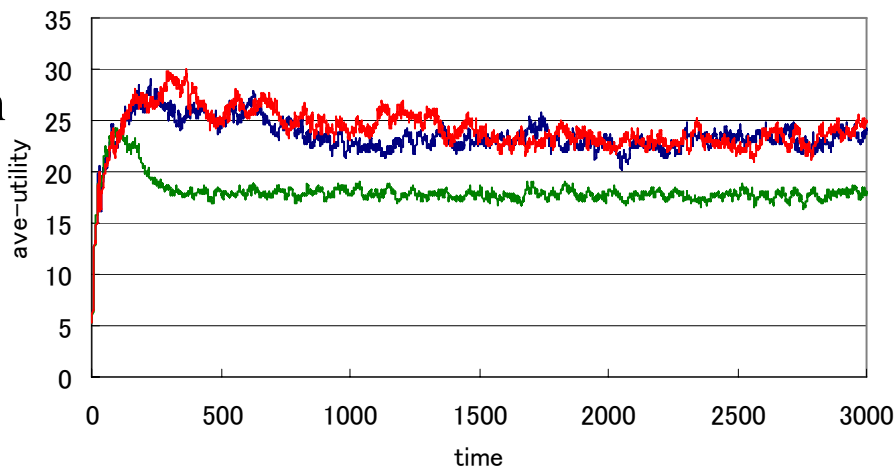
<Model 1> Bilateral trading

<Model 2> Global pricing model

<Model 3> Bilateral trading  
with two-sided matching

The performance of bilateral trading with two-sided matching is very close that of the global market model.

<Utility per agent>





# Conclusion

## <Self-interested Hypothesis vs. Human Sociality Hypothesis>

- : How far agents seek their own interest in a competitive environment?**
- : There are overwhelming evidences that support peoples are also motivated by concerns for fairness and reciprocity.**

**Compromise or cutoff behavior, which is individually irrational, improves the welfare of others.**

**< A small model with the precise analysis vs. A large-scale model with approximation >**

**We may need to explore scaling properties of large-scale socio-economic systems to convince researchers who are so reluctant to give up the the self-interested hypothesis.**